# SpComm3D: A Framework for Enabling Sparse Communication in 3D Sparse Kernels 

Nabil Abubaker<br>nabubaker@inf.ethz.ch<br>ETH Zürich<br>Zürich, Switzerland

Torsten Hoefler<br>htor@inf.ethz.ch<br>ETH Zürich<br>Zürich, Switzerland


#### Abstract

Existing 3D algorithms for distributed-memory sparse kernels suffer from limited scalability due to reliance on bulk sparsity-agnostic communication. While easier to use, sparsity-agnostic communication leads to unnecessary bandwidth and memory consumption. We present SpComm3D, a framework for enabling sparsity-aware communication and minimal memory footprint such that no unnecessary data is communicated or stored in memory. SpComm3D performs sparse communication efficiently with minimal or no communication buffers to further reduce memory consumption. SpComm3D detaches the local computation at each processor from the communication, allowing flexibility in choosing the best accelerated version for computation. We build 3D algorithms with SpComm3D for the two important sparse ML kernels: sampled dense-dense matrix multiplication (SDDMM) and Sparse matrixmatrix multiplication (SpMM). Experimental evaluations on up to 1800 processors demonstrate that SpComm3D has superior scalability and outperforms state-of-the-art sparsity-agnostic methods with up to 20x improvement in terms of communication, memory, and runtime of SDDMM and SpMM. The code is available at: https://github.com/nfabubaker/SpComm3D


## 1 INTRODUCTION

Large-scale sparse kernels such as SDDMM and SpMM are the core operations in many scientific computing and machine learning applications. These kernels involve a sparse matrix with large dimensions, usually an incidence matrix for a graph, and two tall-and-skinny dense matrices. In scientific computing, SpMM is encountered in iterative solvers, when there are multiple right-hand sides, such as block conjugate gradient [27], or blocked eigenvalue algorithms, such as block Lanczos [18] and block Arnoldi [28]. In machine learning and data science, both SpMM and SDDMM became popular for their role in methods used to solve low-rank matrix factorization problems used for recommender systems [24] such as stochastic gradient descent and alternating least squares. A recent survey by Besta et al. [10] demonstrated that SpMM and SDDMM are the backbone of all variants of Graph Neural Networks (GNNs), including Convolutional GNNs and Attentional GNNs. Therefore, improving the parallel scalability of these kernels is pivotal for the feasibility and success of the underlying application.

Executing SDDMM and SpMM, as well as similar kernels, in parallel has been heavily studied and improved. The popular 2D [14, 32] and 3D [3, 25, 30] parallelizations of Dense Matrix-Matrix Multiply (GEMM) provide superior scalability through reduced communication volume and message counts compared to the 1D parallelization. These algorithms inspired the design of parallel algorithms for SpGEMM [13], SDDMM [11], and SpMM [29, 31] with high
success. 2D and 3D algorithms provide nice upper-bounds on the communication bandwidth and latency by dividing the matrix in a structured way leading to well-defined dependency relations that could define the required communication.

Communication in parallel algorithms for sparse kernels is categorized into two categories: sparsity-agnostic bulk communication and sparsity-aware communication. The former relies on communicating data in bulk, without prior knowledge if the receiving processor does or does not require parts of the data communicated. This approach's advantages include easier implementation and utilizing existing efficient algorithms for collective operations (e.g., MPI's Broadcast, All-Reduce, and All-Gather). On the other hand, it usually involves communicating unnecessary data that is not used by the receiving processor, especially when the sparsity is very high. The cost of communicating unnecessary data not only affects the communication bandwidth by adding extra volume, but also requires extra memory for storing this excess data. We empirically demonstrate that bulk communication becomes too expensive with very large sparse matrices in terms volume and memory overheads. Figure 1 shows a simple instance of sparsity-agnostic 2D SDDMM and how it stores and communicates data unnecessary for the computation.

In this work, we present a novel framework that addresses the scalability challenges of distributed-memory sparse kernels by performing sparsity-aware communication while harnessing the power of 2 D and 3 D distributions. We take advantage of the relatively low, and regular, number of messages in 2D and 3D algorithms to address the latency side, and we perform the bare-minimum required communication to address the bandwidth side of the communication overhead. The sparsity-aware communication also reduces memory overhead as it enables the storage of only the necessary data required for the computation, thereby improving scalability on large HPC systems. While there exists 1D [2, 6], "1.5D" [26] (communication-avoiding version of 1D), and 2D [21] algorithms in the literature that utilize sparsity-aware communication, to our knowledge, there exists no 3D algorithm that considers sparsityaware communication.

Our framework fits any 3D sparse kernel that have a computation phase preceded or followed, or both, by a communication phase. SpComm3D does not change the communication based on the sparsity pattern of the input matrix. In this paper we focus on SDDMM and SpMM and we show how to build the sparsityaware 3D algorithms for these kernels with SpComm3D. Our major contributions are summarized as follows:


Figure 1: Unnecessary storage and communication in 2D sparsity-agnostic SDDMM. (a) shows a sample 2D SDDMM on a $2 \times 2$ processor grid. (b) shows the local view of the sub-problem at each processor; data stored locally but unrequired in computation is marked. (c) shows the communication graph and the unrequired communicated data is marked.

- We design a framework that provide a general environment for building large-scale sparse kernels with 2D and 3D distributions that perform sparse communication and require minimal memory footprint (§5).
- Within the framework, we provide several options to perform the sparsity-aware communication revolving around enabling true zero-copy communication in MPI to further reduce memory footprint.
- We outline the communication and memory inefficiencies of the existing 2D and 3D algorithms for SDDMM and SpMM, and we carefully define the minimum required communication to be performed for the correctness of the kernel (§ 4).
- We utilize the new framework for building efficient sparsityaware 3D SDDMM and SpMM algorithms, and we provide a bottom-up guide to doing so in order to pave the way to build other kernels.


## 2 RELATED WORK

2D and 3D algorithms impose upper bounds on communication volume in parallel GEMM operation, and on the number of messages in parallel SpGEMM operation, compared to 1D algorithms. 3D algorithms, such as the work of Agarwal et al. [3], are considered the communication-avoiding version of 2D algorithms as they create multiple copies of the input data to reduce communication. "2.5D" algorithms, first introduced by Solomonik and Demmel [30] and later improved by Kwasniewski et al. [25] are similar to 3D algorithms but control the number of copies created according to available memory. These classes of algorithms are later adapted to SpGEMM: 2D algorithm in the work by Buluç and Gilbert [12, 13], 3D algorithm by Ballard et. al [9], and 2.5D algorithm by Azad et al. [7]. The work by Koanantakool et al. [23] introduced a new class of algorithms called "1.5D" (i.e., the communication avoiding version of 1D) for the sparse-dense matrix-matrix multiplication $\left(\mathrm{SpDM}^{3}\right)$ operation by making redundant copies of the input matrices on top of 1D distribution. The communication in all of these algorithms is sparsity-agnostic.

Several works provided 2D and 3D sparsity-agnostic algorithms for SpMM and SDDMM. Tripathy et al. [31] provided and thoroughly analyzed and compared 1D, 1.5D, 2D, and 3D algorithms for SpMM. Bharadwaj et al. [11] showed how to convert SpMM algorithms to SDDMM, and provided several 2.5D algorithms for SpMM, SDDMM, and FusedMM, a term they coined for the cascade of SDDMM into SpMM, which appears in GNN training and inference. Kannan et al. [19] implemented 2D SpMM within the scope of distributed-memory non-negative matrix factorization algorithm called MPI-FAUN. Selvitopi et al. [29] provides thorough analysis and efficient RDMA-based implementations of several configurations of 2D SpMM.

Sparsity-aware communication, or communication that follows the sparsity of the input matrix, has been previously implemented and analyzed in the context of several 1D algorithms, prominently present in works that utilize graph/hypergraph partitioning for reducing communication overhead in kernels such as sparse matrix-vector multiplication (SpMV) [15, 17, 20, 22], SpMM [2] and SpGEMM [4-6]. Kaya et al. [21] implemented a 2D sparsityaware SpMM algorithm for non-negative matrix factorization on distributed-memory systems. Abubaker et al. [1] implemented a 1D SDDMM algorithm for asynchronous for asynchronous SGD used in matrix factorization. In a recent technical report, Mukhopadhyay et al. [26] utilize sparsity-aware communication for 1.5D SpMM on GPUs. To our knowledge, our work is the first to implement sparsity-aware communication in $2.5 \mathrm{D} / 3 \mathrm{D}$ algorithms.

## 3 PRELIMINARIES

### 3.1 Notations

SpComm3D builds a logical 2D or 3D Cartesian processor grid and distributes the input sparse matrix/matrices onto this grid in a structured fashion. The number of processors in denoted by $P$, the set of all processor by $\mathcal{P}$, and an arbitrary processor in $\mathcal{P}$ by $p_{\alpha}$. For 2D grids, we use the notation $P_{x, y}$ to indicate the processor at location $(x, y)$ of the grid. Similarly, for 3D grids, $P_{x, y, z}$ is located at $(x, y, z)$ in the grid. We use the Matlab notation to address a whole row block, column block, or a slice of the processor grid (e.g., $P_{x, ;, z}$
means the $x$ th row block $\left\{P_{x, 1, z}, P_{x, 2, z}, \ldots, P_{x, Y, z}\right\}$ of the 3D grid and $P_{i,:, z}$ means the $z$ th slice ( 2 D grid) of the 3 D grid).
$\mathbf{S}$ and $\mathbf{C}$ always refer to sparse matrices, whereas $\mathbf{A}$ and $\mathbf{B}$ always refer to dense tall-and-skinny matrices. $\mathbf{S}(i,:$ ) and $\mathbf{S}(:, j)$ denotes the $i$ th row and $j$ th column of $\mathbf{S}$, respectively, and a nonzero element by $s_{i j}$. The $i$ th row in $\mathbf{A}$ or $\mathbf{B}$ is denoted by $\mathbf{a}_{i}, \mathbf{b}_{i}$. The functions $n n z(\cdot)$, nrows $(\cdot)$ and $n \operatorname{cols}(\cdot)$ are respectively used to denote the number of nonzero elements, number of rows, and number of columns of a (sub)matrix. We use $X, Y$ and $Z$ as dimensions for the 2D or 3D processor grid, and $M, N$, and $K$ are used as dimensions for the matrices. $K \ll(M, N)$ is always the number of columns for the tall-and-skinny matrices.

### 3.2 SDDMM and SpMM

Sampled Dense-Dense Matrix Multiplication (SDDMM) $\mathbf{C}=\mathbf{S} \odot \mathbf{A B}^{T}$, where $\odot$ means element-wise multiplication, involves four matrices: $\mathbf{S}$ and $\mathbf{C} \in \mathbb{R}^{M \times N}$ are respectively the sparse sampling and output matrices. $\mathbf{A} \in \mathbb{R}^{M \times K}$ and $\mathbf{B} \in \mathbb{R}^{N \times K}$ are the input dense tall-and-skinny matrices. For large $M$ and $N$ values, materializing the $\mathbf{A B}^{T}$ product becomes prohibitive. Since it will be sampled by the sparse $\mathbf{S}$ matrix, most of the elements in this product are not required, and thus can be computer more efficiently element-by-element as follows:

$$
\begin{equation*}
c_{i j}=s_{i j} \times\left\langle\mathbf{a}_{i}, \mathbf{b}_{j}\right\rangle, \forall(i, j) \in\left\{(i, j) \mid s_{i j} \neq 0\right\} \tag{1}
\end{equation*}
$$

Sparse times Dense Matrix Multiplication (SpMM) A $=\mathbf{S B}^{T}$ involves $\mathbf{S}, \mathbf{B}$ as input, and $\mathbf{A}$ as output, all of sizes described above. SpMM can be computed row-by-row in serial execution as

$$
\begin{equation*}
\mathbf{a}_{i}=\sum_{1 \leq j \leq N} s_{i j} \times \mathbf{b}_{j} \tag{2}
\end{equation*}
$$

Both SDDMM and SpMM operate on one sparse matrix and two dense tall-and-skinny matrices. Both operations require taking/outputting two dense vectors per nonzero element. In SDDMM, both vectors are taken as input, whereas in SpMM $\mathbf{b}_{j}$ is taken as input and $\mathbf{a}_{i}$ is given as output. Therefore, most of the analyses and algorithms explained afterwords apply to both operations. Our presentation, figures, and analyses will focus on SDDMM hereafter, and we will explain the difference if required.

To parallelize (1) and (2), the fine-grain tasks are first identified, then the communication is determined following how these fine-grain tasks are distributed to different processors and the dependency between them. For SDDMM, we define the fine-grain task as computing one scaled inner product. Each fine-grain task requires two vectors (each of size $K$ words) and one scalar. For SpMM, we define the fine-grain task as computing one vector scaling as in (2). Each fine-grain task requires one vector of size $K$ words and a scalar.

### 3.3 Sparsity-agnostic parallel algorithms for SDDMM and SpMM

Parallel SDDMM and SpMM algorithms are categorized according to how the input sparse matrix is partitioned among processors. The partitioning can either be based on granularity or on structure.

Granularity-based categorization can be either fine-grain (nonzerobased) or coarse-grain (row or column-based). Structure-based categorization relies on the dividing the computational iteration space in a structured manner into $1 \mathrm{D}, 2 \mathrm{D}$, or 3 D shape, which mirrors a partitioning on the dimensions of the matrix itself. We provide an overview of such algorithms and provide a volume upper bound for the best algorithm in each category. Figure 2 illustrates 1D, 2D, and 3D algorithms for SDDMM. Since our concrete goal is to reduce the communication volume via sparsity-aware communication, we analyze the communication of the sparsity-agnostic algorithms in terms of the received volume per processor.

In 1D algorithms, the sparse matrix $\mathbf{S}$ is divided into $P$ parts either row-wise or column-wise, necessitating the communication of one of the dense matrices, not both. Assuming 1D division on rows, the rows of the dense matrix $\mathbf{A}$ are also divided conformably with the rows of $\mathbf{S}$. Naturally, a processor $p_{\alpha}$ is assigned the ownership of the dense rows of $\mathbf{A}$ that align with its assigned row block of $\mathbf{S}$, which necessitates no communication on the rows of $\mathbf{A}$. On the other hand, the dense matrix $\mathbf{B}$ is likely to be used by all processors thus necessitates communication. Assuming $\mathbf{B}$ is divided into $P$ parts, and each part is assigned to a single processor, each processor will communicate its part to $P-1$ processors, leading to approximately

$$
\mathbf{B}_{\text {size }} \times \frac{P-1}{P}
$$

words of volume to be received per processor. The per-processor memory requirement for storing the dense matrices is

$$
\frac{\mathbf{A}_{\text {size }}}{P}+\mathbf{B}_{\text {size }}
$$

In 2D algorithms, the sparse matrix is divided into a $\sqrt{P} \times \sqrt{P} 2 \mathrm{D}$ grid. Depending on how the dense matrices are divided, different versions of the 2D algorithms emerged [29]. We consider the version where $\mathbf{A}$ and $\mathbf{B}$ are 1D divided row-wise into $P$ parts, which is wellsuited for single-step SpMM and SDDMM algorithms [11, 19, 29].

Each of the inner $\sqrt{P}$ parts of $\mathbf{A}$ along a row block of processors $P_{x,:}$ is assigned ownership to one of the processors in that block, similarly fo $\mathbf{B}$. The 2D grid partitioning restricts the requirement of a given dense matrix part to only $\sqrt{P}$ processors along the same row or column block. The communication in 2D algorithms is required along both dimensions. Each processor receives an approximate of

$$
\left(\mathbf{A}_{s i z e}+\mathbf{B}_{s i z e}\right) \times \frac{\sqrt{P}-1}{P}
$$

words of volume. The per-processor memory requirement for storing the dense matrices is

$$
\frac{\mathbf{A}_{\text {size }}+\mathbf{B}_{\text {size }}}{\sqrt{P}}
$$

3D algorithms, sometimes called 2.5D algorithms [11, 31], are considered the communication-avoiding version of the 2D algorithms. These algorithms add an additional dimension $Z$ to the 2D algorithms by running $Z$ instances of the sparse kernel concurrently, where each instance is responsible for computing $K / Z$ columns of the dense matrix. This naturally means that the each instance is resembled with a replica of the sparse matrix partitioned according to the 2D scheme into $\sqrt{P / Z} \times \sqrt{P / Z}$ grid. The columns of the dense matrix are partitioned into $Z$ parts, each part is used


Figure 2: Sparsity-agnostic 1D, 2D, and 3D SDDMM. The communication required stays the same whether $S$ is dense or sparse.
exclusively by a separate instance. With this scheme, each processor receives an approximate of

$$
\left(\mathbf{A}_{s i z e}+\mathbf{B}_{s i z e}\right) \times \frac{\sqrt{\frac{P}{Z}}-1}{P}
$$

words of volume. See Figure 2 (right). The per-processor memory requirement for storing the dense matrices is

$$
\frac{\mathbf{A}_{\text {size }}+\mathbf{B}_{\text {size }}}{Z \sqrt{\frac{P}{Z}}} .
$$

## 4 SPARSITY-AWARE COMMUNICATION ANALYSIS OF 3D SDDMM

In a row processor block $P_{x, ;, z}$, let $\lambda_{i}$ be the number of processors at which row $\mathbf{S}(i,:)$ has at least one nonzero element, and $\Lambda_{i}$ be the set of such processors. Assuming that row $\mathbf{a}_{i}$ of dense matrix A is owned by one of the processors in $\Lambda_{i}$, it is required by $\lambda_{i}-$ 1 processors. A similar discussion holds for a column processor block $P_{i, y, z}$ and the dense matrix B. Then, the total communication volume exchanged in sparsity-aware SDDMM is equal to

$$
\sum_{i \in[[1, \operatorname{rrows}(\mathbf{S})]]}\left(\lambda_{i}-1\right)+\sum_{j \in[[1, n \operatorname{cols}(\mathbf{S})]]}\left(\lambda_{j}-1\right)
$$

In the dense case, $\lambda_{i}$ and $\lambda_{j}$ become equal to sizes of the $X$ and $Y$ dimensions of the 3D $\operatorname{grid}\left(\sqrt{\frac{P}{Z}}\right)$ which leads to correct total volume.

The communication requirement modelled by $\lambda$ depends on three factors: (i) the sparsity pattern of $\mathbf{S}$, (ii) the number of processors and (iii) how the matrix nonzero elements are distributed to processors. While factor (i) is vital, we assume the sparsity pattern of $\mathbf{S}$ is irregular (i.e., non-uniformly distributed) in order for (iii) to take effect. With the sparsity-agnostic communication, a row/column of $\mathbf{S}$ is divided into $P$ parts incurs $K(P-1)$ words of communication volume. This means that the communication volume scales with how many parts a row/column partitioned into. With $\lambda$, on the other hand, the volume depends on how many of the $P$ parts has at least one nonzero element. As $P$ grows larger, the probability of a certain part of a row/column of $\boldsymbol{S}$ is empty becomes higher. Therefore, the communication requirement according to $\lambda$ is loosely related to the growing $P$.


Figure 3: A sample matrix distributed onto $5 \times 5$ grid and the $\lambda$ values of some rows/columns.

In order to define the received volume per processor, we define two sets $I_{\alpha}=\left\{i \mid \mathbf{S}_{\alpha}(i,:) \neq \emptyset \wedge \operatorname{owner}\left(\mathbf{a}_{i}\right) \neq p_{\alpha}\right\}$ and $\mathcal{J}_{\alpha}=$ $\left\{j \mid \mathbf{S}_{\alpha}(:, j) \neq \emptyset \wedge\right.$ owner $\left.\left(\mathbf{b}_{j}\right) \neq p_{\alpha}\right\}$ as respectively the sets of row and column indices such that $p_{\alpha}$ has at least one zero in the row/column of $\mathbf{S}$ in those indices and $p_{\alpha}$ is not the owner of the corresponding dense row. Then, the volume received by $p_{\alpha}$ is equal to $\frac{K}{Z}\left(\left|\mathcal{I}_{\alpha}\right|+\left|\mathcal{J}_{\alpha}\right|\right)$ words of volume.

Our goal in this work is to achieve the $\lambda$-based communication requirement defined in this section using efficient communication without extra unnecessary volume. Building the structure of this communication for most sparse kernels can be cumbersome, which is why we introduce SpComm3D to provide a structured and convenient environment for building sparse kernels with minimum communication.

## 5 THE SPCOMM3D FRAMEWORK

The framework has three design goals: (i) communicate only required data, (ii) minimal memory footprint, and (iii) communicationagnostic local computation at each processor to allow utilizing existing efficient algorithms/tools for sparse/dense kernels.

### 5.1 General structure and assumptions

SpComm3D assumes that the sparse kernel is used in a larger context, and is repeated multiple times. The sparsity pattern of the input sparse matrix/matrices remains fixed, while the values might
be updated. SpComm3D also works under the assumption that there are other dense or sparse data that will be computed or used during the sparse kernel, and this data is updated at each iteration either before or after executing the sparse kernel.

SpComm3D revolves around a local computation, and communicates/stores the minimum amount of data required for the overall correctness of this computation. Therefore, computing a sparse kernel with SpComm3D naturally involves three phases: PreComm, Compute, and PostComm. PreComm gathers the data required for computation from different processors, while PostComm communicates partial results to the processors responsible for holding the final value.

Compute is the local computation phase at each processor. The assumption here is that the local computation is agnostic to the communication and the general structure of the problem, thus should be operating on local data with localized indices. The PreComm phase ensures that the data used in the Compute phase is correct and up-to-date, bridging the gap between the global and local views of the problem.

Data communicated in PreComm and PostComm is represented in local memory by a data segment, which may consist of one or more data words and is identified globally using a unique ID typically related to the sparse matrix's global row/column indices. We call such data segment Data Unit (DU) hereafter. A DU might be required or updated by several processors but owned by only one, which can be retrieved with owner $(\cdot)$.

As per our assumption that the sparse kernel will be executed multiple times with the same sparse matrix, we minimize the amount of work to be done during the PreComm and PostComm phases by introducing a setup phase that is executed once. This phase builds all the communication structures, buffers, and meta information that are used in the communication. Then, the PreComm and PostComm phases are used to merely move data between processors. A similar philosophy is followed in MPI's persistent communication.

### 5.2 Data distribution

We assume that the input matrices $\mathbf{S}$ is distributed onto the $2 \mathrm{D} / 3 \mathrm{D}$ processor grid as follows: the matrix is partitioned into $X \times Y$ in the row/column dimension spaces. Each block $\mathbf{S}_{x, y}$ of the partitioned matrix is assigned to processor $P_{x, y}$ in a 2D grid, and then partitioned into $Z$ parts in the nonzero space. That is, $\mathbf{S}_{x, y}$ is distributed equally among $Z$ processors in $P_{x, y, \text { : }}$ as $\mathbf{S}_{x, y}^{1}, \cdots, \mathbf{S}_{x, y}^{Z}$ in a 3D grid. We refer to this distribution as Dist2D or Dist3D depending on the grid used.

Each matrix block $\mathbf{S}_{x, y}^{z}$ is localized to its processor by localizing row/column indices and removing empty rows/columns. In order to keep the global information of the local sparse matrix, SpComm3D stores two arrays along the rows and the columns: globalMap and localMap. globalMap stores the global index of each local row/column and localMap stores local indices of the rows/columns in the sub-matrix.

Apart from the local sparse matrix, each processor might hold additional local data assigned to it. This data can be another sparse matrix, a dense matrix, or a vector of elements. The local data can be abstracted as DUs.


Figure 4: Local SDDMM view of processor $p_{3}$ (cf. Figure 1) after applying SpComm3D localization. The global and local indices are mapped to each other with globalMap and localMap.

### 5.3 Sparse communication model

Sparse communication can be very irregular, and is best defined with a point-to-point ( P 2 P ) communication graph commG $=$ $(\mathcal{P}, \mathcal{M})$. While $\mathcal{P}$ is the set of all processors as defined previously, $\mathcal{M}$ is the set of all messages between these processors. A message $m_{\alpha \rightarrow \beta} \in \mathcal{M}$ exists if there is at least one DU that either owned by $p_{\alpha}$ and required by $p_{\beta}$, or partially computed by $p_{\alpha}$ and owned by $p_{\beta}$.

During PreComm, the messages typically constitute DUs that are owned by one processor and required by others (broadcast). During this phase, a DU can appear in several outgoing messages, whereas incoming messages contain unique DUs. This is because each DU is owned by a single processor, and it is not possible to receive it from multiple processors.

During PostComm, the messages are typically partial results of DUs that will be sent to their respective owners (reduce). During this phase, a DU can appear in several incoming messages, whereas outgoing messages contain unique DUs.

The distinction whether incoming/outgoing messages contain unique, or otherwise non-unique, DUs is important for the discussion on how to we enable true zero-copy sparse communication in SpComm3D.

The messages are exchanged using either point-to-point communication (MPI_Send/MPI_Recv or their non-blocking equivalents). We propose to exchange these messages without relying on send/receive buffers. For the sake of completeness, we propose three versions of handling buffers in sparse communication where the first method assumes the use of send/receive buffers.
5.3.1 Method1: Sparse Communication with Both Buffers (SpC-BB). The straight forward way to build the messages in comm $G$, given that global IDs to be in the message are known, is to go over each global ID and copy its associated DU to a buffer. Similarly, at the receiver's end, copy the received DU one by one to their correct location in memory according to their global ID. This approach adds extra costs when compared to the sparsity-agnostic bulk communication approach, which are the cost of extra memory resembled in the send and receive buffers, as well as the cost of data movement while copying to/from the local data. This method works regardless of the uniqueness of DUs in outgoing/incoming messages.


Figure 5: Zero-copy techniques based on DU uniqueness within incoming/outgoing message.
5.3.2 Method2: Sparse Communication with Send/Recv Buffer Only (SpC-SB/SpC-RB). Our first improvement over SpC-BB it to get rid of either the send or the receive buffers by re-arranging how the local data is stored to align with the order of the sent/received data. In other words, assume processor $p_{\beta}$ sends/receives from $n$ processors, $1 \leq n \leq Y$, in $P_{x, ;, z}$ in the following sequence $\left\langle p_{i_{1}}, p_{i_{2}}, \ldots, p_{i_{n}}\right\rangle$ and the set of DUs received by $p_{i_{n}}$ is denoted by $\mathcal{D}_{i_{n}}$. Then, the local data at the processor is stored as

$$
\left[\mathcal{D}_{i_{1}}, \mathcal{D}_{i_{2}}, \ldots, \mathcal{D}_{i_{n}}\right]
$$

and the messages are directly sent/received at the position starting from the first index of $\mathcal{D}_{i_{1}}$.

With this method, it is only possible to remove the buffers for either the outgoing or incoming messages, whichever has unique DUs. This is because, if the messages have non-unique DUs the relation between the duplicates of the DUs in several messages and their unique location in local memory becomes many-to-one rather than one-to-one, hence making the local data alignment impossible.
5.3.3 Method3: $\boldsymbol{S p a r s e}$ Communication with $\mathbf{N o}$ Buffers (SpC-NB). For messages with non-unique DUs, we utilize MPI_Type_Indexed to address the multiple copies of the same DU without the need to explicitly moving them to a buffer. MPI_Type_Indexed creates a new datatype by aligning blocks of pre-defined datatype. The MPI_Type_Indexed type is initialized with the starting address of a contiguous memory chunk, and each block is defined with length and displacement from the starting address. It is possible to create two (or more) copies/projections of the same local DU in two messages $M_{\alpha \rightarrow \beta}$ and $M_{\alpha \rightarrow \gamma}$ (or more) by simply adding the DU as a block in the MPI_Type_Indexed object of each message with the same displacement.

We consider a DU to be the minimum block in MPI_Type_Indexed. If multiple DUs appear consecutively in memory, we merge them into a single block in order to minimize the data required to build the MPI_Type_Indexed object.

SpComm3D takes the communication graph from the user in the init phase. Based on this graph, the SpComm3D framework internally handles all the steps necessary to perform efficient communication during the PreComm phase, including the zero-copy implementations. SpComm3D returns a communication object per communication graph in the setup phase. Assuming the PreComm communication object is called precomm, then during the PreComm phase the user simply calls precomm. communicate() to perform the efficient sparse communication.

## 6 BUILDING 3D SDDMM AND SPMM ALGORITHMS WITH SPCOMM3D

We begin by showing in detail the steps to build 3D SDDMM algorithm with SpComm3D, and later we reuse some of these steps for 3D SpMM. Using Dist3D, we can formulate the 3D parallel SDDMM algorithm as follows: the $\mathbf{A}$ and $\mathbf{B}$ matrices are partitioned row-wise into $X$ and $Y$ parts conformably with the rows and columns of $\mathbf{S}$, respectively. Then, $\mathbf{A}$ and $\mathbf{B}$ are partitioned column-wise to $Z$ parts $\mathbf{A}^{1}, \cdots, \mathbf{A}^{z}, \cdots, \mathbf{A}^{Z}$. Here, a part $\mathbf{A}_{x}^{z}$ is used by processors in $P_{x,,, z}$. Similarly, a part $\mathbf{B}_{y}^{z}$ is used by processors in $P_{i, y, z}$ We assume each processor in $P_{x, ; z}$ owns almost equal number of rows of $\mathbf{A}_{x}^{z}$. This naturally means that $\mathbf{A}_{x}^{z}$ is further divided into $Y$ parts. We abuse the 2D notation of matrix blocks to refer to 1D internal parts of a dense 1D matrix block. For instance, $\mathbf{A}_{x, y}^{z}$ refers to the 1D block $\mathbf{A}_{x X+y}^{z}$. The processor that owns a row $\mathbf{a}_{i}^{z}$ is retrieved with the owner $\left(\mathbf{a}_{i}^{z}\right)$.

We start by defining the requirements of the Compute phase, and then move to explain how to fulfill these requirements in the init, PreComm, and PostComm phases.

### 6.1 The Compute phase

In 3D SDDMM, the responsibility of a processor $P_{x, y, z}$ is to compute partial results of inner products for all nonzeros in $\mathbf{C}_{x, y}$, and then be responsible for holding the final results of nonzeros in $\mathbf{C}_{x, y}^{z}$. Here, the two vectors in the inner product $\left\langle\mathbf{a}_{i}, \mathbf{b}_{j}\right\rangle$ that take place for each nonzero element in $\mathbf{S}_{x, y}$ are of size $\frac{K}{Z}$. This clearly resembles the shift from 2D to 3D algorithms as dividing the fine-grain tasks of computing the inner product of two dense vectors of size $K$ into $Z$ sub-tasks each of size $\frac{K}{Z}$.

Initially, $P_{x, y, z}$ owns $\mathbf{S}_{x, y}^{z}$ only. Therefore, it requires $\mathbf{S}_{x, y} \backslash \mathbf{S}_{x, y}^{z}$ as well as all (sub)rows of $\mathbf{A}$ and $\mathbf{B}$ required for the computation of $\mathbf{C}_{x, y}$. With the assumption that the values and sparsity patter of $\mathbf{S}_{x, y}$ are constant during multiple iterations of computing SDDMM, the gathering of $\mathbf{S}_{x, y}$ can be done at the init phase. On the other hand, the values in $\mathbf{A}$ and $\mathbf{B}$ will change at each iteration, therefore they should be communicated at the PreComm phase.

The local compute phase at each processor is stand-alone by the design of the framework. A processor $P_{x, y, z}$ has a localized version of $\mathbf{S}_{x, y}$ as well as all required $\mathbf{A}$ and $\mathbf{B}$ rows by the end of the PreComm phase. Figure 4 shows an example of the state of local SDDMM. The Compute phase is performed following (1), either on CPU, GPU or any accelerator for sparse computations using state-of-the-art sequential or shared-memory-parallel codes.

### 6.2 The PreComm phase

In this phase, the dense $\mathbf{A}$ - and $\mathbf{B}$-matrix rows are communicated. Computing partial $\mathbf{C}_{x, y}$ by $P_{x, y, z}$ requires dense rows from $\mathbf{A}_{x}^{z}$ and $\mathbf{B}_{y}^{z}$. Since $\mathbf{A}_{x}^{z}$ is required by all processors in $P_{x, ;, z}$, we assume that $P_{x, y, z}$ owns an equal part of $\mathbf{A}_{x}^{z}$. Similar discussion hold for $\mathbf{B}_{y}^{z}$.

Since the core goal of SpComm3D is to perform sparse communication, we follow the $\lambda$-based communication discussed in Section 4. A message $m_{\alpha \rightarrow \beta}$ from processor $p_{\alpha}$ to processor $p_{\beta}$, where $p_{\alpha}, p_{\beta} \in P_{x, ;, z}$ is formed as

$$
\begin{equation*}
m_{\alpha \rightarrow \beta}=\left\{\mathbf{a}_{i} \mid p_{\alpha}, p_{\beta} \in \Lambda_{i} \wedge \operatorname{owner}\left(\mathbf{a}_{i}\right)=p_{\alpha}\right\} \tag{3}
\end{equation*}
$$

Similarly, a message $m_{\gamma \rightarrow \delta}$ from processor $p_{\gamma}$ to processor $p_{\delta}$, where $p_{\gamma}, p_{\delta} \in P_{i ; y, z}$ is formed as

$$
\begin{equation*}
m_{\gamma \rightarrow \delta}=\left\{\mathbf{b}_{j} \mid p_{\gamma}, p_{\delta} \in \Lambda_{j} \wedge \operatorname{owner}\left(\mathbf{b}_{j}\right)=p_{\gamma}\right\} \tag{4}
\end{equation*}
$$

### 6.3 The PostComm phase

Gathering the partial results of $\mathbf{C}_{x, y}$ requires that processor $P_{x, y, z}$ receives all partial results of the nonzero elements it owns from the other $Z-1$ processors in $P_{x, y,:}$. This amounts to receiving $(Z-1) \times n n z\left(\mathbf{S}_{x, y, z}\right)$ words of data. With highly sparse matrices, the cost of this phase is very low compared to PreComm. We perform this operation with a Reduce-Scatter rather that converting it to sparse communication primitives.

### 6.4 The Setup phase

All the configurations required for the PreComm, Compute, and PostComm phases are performed in this phase. These configurations are performed once and used multiple times. The first is gathering $\mathbf{S}_{x, y}$ at each processor in $P_{x, y,:}$. This is done with an All-Gather operation on $\mathbf{S}_{x, y}^{z}$ by all processors in $P_{x, y,:}$.

The second configuration is distributing the dense A and $\mathbf{B}$ among processors. In Section 6.2 we stated our assumption that the dense rows are distributed equally (owned) to the processors that use them. However, without careful attention, a random equal distribution will invalidate the $\lambda$-based communication discussed in Section 4.

The $\lambda$-based communication of a dense row $\mathbf{a}_{i}$ is accurate only if the processor that owns $\mathbf{a}_{i}$ is part of $\Lambda$. Otherwise, a dense row $\mathbf{a}_{i}$ that is assigned to a processor outside of $\Lambda_{i}$ will incur an extra unnecessary communication and storage of size $K$ words.

We propose Algorithm 1 to efficiently distribute the dense rows to processors in their respective $\Lambda$ sets in parallel. The algorithm starts by dividing the work to be done by each processor. Each processor is responsible for finding an owner of a set of rows, and these rows can be assigned randomly in a load balanced manner (lines 1-2). Then, each processor loops over the rows it uses and send their global index to the processor responsible for their assignment (lines 4-13). After each processor receives the list of candidates to each row, it loops over the candidates and picks one processor at random to assign it as the owner of the corresponding row (lines 14-22). The candidates that each processor receives for each row index are the processors that actually use that index in their local computation, i.e., they have at least one nonzero element with such an index. For this reason, the processor that is picked at random here will certainly be one of the processors involved in the communication of the target row, thus not incurring any additional

```
Algorithm 1 Parallel \(\lambda\)-aware random distribution on \(p_{\alpha}\)
    Input \(\operatorname{nrows}(\mathbf{S})\) or \(\operatorname{ncols}(\mathbf{S})\) as gsize, \(\operatorname{nrows}\left(\mathbf{S}_{\alpha}\right)\) or
nrows \(\left(\mathbf{S}_{\alpha}\right)\) as dsize, P, globalMap
    Output owner array owner
Require:
    myrows \(\leftarrow\) the myrank chunk of dsize/P
    candidates \(\leftarrow\) array of lists of size \(\mid\) myrows \(\mid\)
    sendInfo \(\leftarrow\) array of lists of size \(P\)
    for \(i\) in range ( 1, ldsize) do
        \(g i \leftarrow\) globalMap \([i]\)
        \(p \leftarrow\) processor responsible for assigning \(g i\)
        if \(p=\) myrank then
            candidate \([g i] \leftarrow\) candidates \([g i] \cup p\)
        else
            sendInfo \([p] \leftarrow \operatorname{sendInfo}[p] \cup g i\)
        end if
    end for
    recvInfo \(\leftarrow\) Exchange sendInfo
    for \(p\) in range \((1, P)\) do
        for row ID rid in recvInfo \([p]\) do
                candidates \([\) rid \(] \leftarrow\) candidates \([\) rid \(] \cup p\)
        end for
    end for
    myowner \(\leftarrow\) empty array of size |myrows \(\mid\)
    for row ID rid in myrows do
        myowners \([\) rid \(] \leftarrow\) pick a random \(p \in\) candidates \([r i d]\)
    end for
    All-Gather(myOwner, owner)
```

unnecessary communication. Finally, an All-Gather operation is performed in order to gather the ownership information from all processors (line 23).

The third configuration that will be used multiple times during the parallel SDDMM is building the communication graph of the PreComm phase. Based on this graph, the SpComm3D framework internally handles all the steps necessary to perform efficient communication during the PreComm phase, including the zero-copy implementations.

### 6.5 Building SpMM with SpComm3D

The process to building SpMM is similar to that of SDDMM. The distribution of the sparse matrix is the same as that of SDDMM. In the Compute phase, a processor $P_{x, y, z}$ is responsible for computing partial results of dense rows in $A_{x, y}^{z}$. For that, $P_{x, y, z}$ also needs the whole $\mathbf{S}_{x, y}$, which is gathered during the init phase. The SpMM requires B-matrix rows in the PreComm phase. The communication graph is built with (4). After the Compute phase, $P_{x, y, z}$ sends partial results for all the dense rows that it does not own to their respective owners during the PostComm phase. The communication graph in this phase is constructed with (3) but replacing owner $\left(\mathbf{a}_{i}\right)=p_{\alpha}$ by owner $\left(\mathbf{a}_{i}\right)=p_{\beta}$. This is because the owner is the receiving party not the sending party as this is a reduce operation. The communication cost of the PreComm phase of SDDMM can be thought of as distributed among both PreComm and PostComm phases in SpMM.

In other words, both PreComm and PostComm phases in SpMM are of equal importance, unlike SDDMM.

## 7 EXPERIMENTAL EVALUATION

Our evaluation goal is to empirically assess the theoretical aspects of SpComm3D in terms of reducing communication and memory footprint, and how they reflect on actual runtime and scalability, compared to sparsity-agnostic 3D algorithms.

In line with the presentation of the paper, we consider parallel SDDMM and SpMM in our experiments. We use real-world sparse matrices with at least 100 M nonzero elements, and we vary three variables, $K, Z$, and $P$ to study their effect on both SpComm3D and the sparsity-agnostic 3D algorithms. Unless stated otherwise, all runtime results reported in this section are result of averaging five different runs per experiment.

For the sparsity-agnostic 3D algorithm (§3.3), we provide our own implementation (referred to as Dense3D hereafter), and also compare against the state-of-the-art framework Half-and-Half $(\mathrm{HnH})$ by Bharadwaj et al. [11] that provides the same algorithm under the name "2.5D sparse replicating" (referred to as HnH hereafter). This existing framework is compared against PETSc [8] and shown to outperform it significantly. For this reason, we refrain from comparing against PETSc. We use the abbreviations in Section 5.3: $\mathrm{SpC}-\mathrm{BB}, \mathrm{SpC}-\mathrm{R} / \mathrm{SB}$, and $\mathrm{SpC}-\mathrm{NB}$ to refer to the specific implementation in SpComm3D, and we use the framework's name when comparing metrics irrelevant to the implementation such as communication volume.

### 7.1 Dataset and Experimental Setting

Table 1: Sparse matrices used in our experiments

| Matrix | \#rows/cols | \#nonzeros | Density |
| :--- | ---: | ---: | ---: |
| arabic-2005 | $22,744,080$ | $639,999,458$ | $1.24 \times 10^{-06}$ |
| delaunay_n24 | $16,777,216$ | $100,663,202$ | $3.58 \times 10^{-07}$ |
| europe_osm | $50,912,018$ | $108,109,320$ | $4.17 \times 10^{-08}$ |
| GAP-kron | $134,217,726$ | $4,223,264,644$ | $2.34 \times 10^{-07}$ |
| GAP-road | $23,947,347$ | $57,708,624$ | $1.01 \times 10^{-07}$ |
| GAP-web | $50,636,151$ | $1,930,292,948$ | $7.53 \times 10^{-07}$ |
| kmer_A2a | $170,728,175$ | $360,585,172$ | $1.24 \times 10^{-08}$ |
| twitter7 | $41,652,230$ | $1,468,365,182$ | $8.46 \times 10^{-07}$ |
| uk-2002 | $18,520,486$ | $298,113,762$ | $8.69 \times 10^{-07}$ |
| webbase-2001 | $118,142,155$ | $1,019,903,190$ | $7.31 \times 10^{-08}$ |

We evaluate with ten real-world sparse matrices that represent graphs, obtained from the The SuiteSparse Matrix Collection ${ }^{1}$ [16]. The properties of the sparse matrices are detailed in Table 1. All our matrices have between 100 million and 4.2 billion nonzero elements.

We implemented SpComm3D as well as Dense3D using C++ and used MPI for inter-process communication. All our experiments are taken on CSCS Piz Daint HPC system based in Switzerland. We use the CPU partition of the Cray XC40/XC50 system, which is equipped with 1813 dual-socket Intel Xeon E5-2695 processors clocked at 2.10 GHz and has 64 GiB of DD3 RAM memory. The

[^0]nodes are connected with a Cray Aries network that uses a dragonfly network topology. All the codes are compiled with a Cray-clang compiler and a system-provided Cray-MPICH on SUSE Linux Enterprise Server 15-SP2 operating system.

### 7.2 High-level total runtime comparison with sparsity-agnostic algorithms

We begin by comparing SpComm3D against Dense3D and HnH in terms of total runtime of five iterations of SDDMM followed by SpMM . The reason for this comparison is that HnH is designed to perform fused SDDMM and SpMM operations, and has restrictions when it comes to choosing $Z$ and $K$ values. Also, the $X$ and $Y$ dimensions of the 3D mesh should be equal in HnH. SpComm3D does not have such restrictions, and works with any $X, Y, Z$, and $K$ values as long as the trivial restrictions hold: $K / Z \geq 1$ and $X \times Y \times Z=P$. For this reason, we set $Z$ and $K$ parameters to values that work for HnH , which are respectively 4 and 60 . We later vary these values in experiments that do not include HnH .


Figure 6: Comparing SpComm3D, Dense3D, and HnH in terms of five SDDMM followed by SpMM operations runtime.

Figure 6 compares SpComm3D, Dense3D, and HnH in terms of runtime of five SDDMM followed by SpMM operations on 900 processors. HnH as is without any modifications by selecting the proper 2.5D algorithm " 2.5 D sparse replicating". As the figure shows, SpComm3D significantly outperforms both Dense3D and HnH. Although theoretically the same algorithm, HnH and Dense3D show different runtimes on different matrices. While most instances show that both methods perform similarly, or slightly in favor of one of them, three instances show significant difference in favor of Dense3D. This behavior could be explained by the fact that HnH uses multiple blocking send/recieve calls (MPI_Sendrecv) per processor to realize the All-Gather operation required for the first phase of communication, whereas Dense3D uses non-blocking broadcasts (MPI_Ibcast) to realize the same communication.


Figure 7: Strong Scaling (log-log) of SDDMM with $K=120$ and $Z=4$. A missing value means infeasible run (out-of-memory).

### 7.3 Strong Scaling

In order to study the scaling of SDDMM with SpComm3D compared to Dense3D, we run SDDMM on 36, 72, 180, 360, 540, 900, and 1800 processors. Within SpComm3D, we compare two implementations of sparsity-aware communication: $\mathrm{SpC}-\mathrm{BB}$ and $\mathrm{SpC}-\mathrm{NB}$. We fixed $K$ value to 120 and $Z$ value to 4 . Figure 7 shows the strong scaling results for all matrices in our dataset. In the figure, there are some missing values that we were unable to obtain due to high memory demand (out-of-memory errors). Most of such missing values are from the Dense3D method, and/or with experiments taken on small number of processors ( $P=180$ or below).

As seen in the figure, sparsity-aware communication is superior to Dense3D in terms of runtime and memory scalability. There is a significant gap between Dense3D and the other sparsity-aware methods in terms of runtime, especially as $P$ becomes higher. When comparing sparsity-aware methods against each other, it is clear that SpC-BB is inferior to SpC-NB in all cases when $P<900$. When $P$ grows larger than $900, \mathrm{SpC}-\mathrm{BB}$ performs similar to $\mathrm{SpC}-\mathrm{NB}$ in six out of ten matrices.

### 7.4 Harnessing Sparsity: evaluation of memory and communication

As demonstrated in this paper, using sparsity-aware communication in SpComm3D not only reduces the actual volume of communication between processors, but also enables reducing the overall memory footprint by avoiding the storage of any unnecessary dense $\mathbf{A}$ or $\mathbf{B}$ rows. We have measured the total memory required for storing dense A and B on 1800 processors when $K$ value is set to 240. We used three matrices from our dataset, and varied the $Z$ dimension between 2,4 , and 9 , leading to three different 3D mesh configurations. Figure 8 shows that the overall memory consumption of SpComm3D is significantly lower than that of Dense. On "arabic-2005", the memory is reduced by 2.5 x to 3.5 x , depending on the value of $Z$. On "kmer_A2a", the memory is reduced by 5 x to 10 x , effectively saving 4 TB to 10 TB of overall memory. Similar ratios appear with "webbase-2001". The memory consumption of Dense3D is reduced with increasing $Z$ value, an expected behavior


Figure 8: Memory, volume, and SDDMM runtimes of Dense3D vs. SpComm3D on three matrices on 1800 processors (with $K=120$ ).
following the theoretical analysis. The memory consumption of the sparsity-aware algorithm also decreases in a slower rate with increasing $Z$ compared to Dense3D.

### 7.5 Overall communication volume and runtime improvements

In order to assess SpComm3D more thoroughly, we report our SDDMM runs using Dense3D, SpC-BB, SpC-RB, and SpC-NB with $Z=\{2,4,9\}$ and $K=\{60,120,240\}$ on 900 and 1800 processors.

Table 2: Comparing Dense3D vs all methods of SpComm3D in terms of max receive volume and total SDDMM runtime on 900 processors with different $Z$ and $K$ values.

| $Z$ | Method | $\begin{aligned} & \text { Max. Recv } \\ & \text { Volume } \\ & \text { (K-normalized) } \end{aligned}$ | SDDMM runtime (ms) |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | $K=60$ | 120 | 240 |
| $Z=2$ | Dense3D | 2,129,152 | 3,021 | 6,063 | 7,798 |
|  | $\overline{\mathrm{Sp}} \overline{\mathrm{C}}-\overline{\mathrm{B}} \cdot \bar{\square}$ |  | $\overline{795}$ | 1,600 | 3,597 ${ }^{-}$ |
|  | SpC-RB | 328,264 | 720 | 1,434 | 2,702 |
|  | SpC-NB |  | 714 | 1,295 | 2,587 |
|  | Improvement | 6.5x | $\overline{4.2} \overline{\mathrm{x}}$ | $\overline{4} .7 \mathbf{x}$ | $\overline{\mathbf{3} .0 \times}{ }^{-}$ |
| $Z=4$ | Dense3D | 1,453,978 | 2,569 | 5,065 | 7,506 |
|  | $\overline{\mathrm{Sp}} \overline{\mathrm{C}}-\overline{\mathrm{B}}$ - |  | $\overline{9} 02$ | 1,773 | 3,772 |
|  | SpC-RB | 289,876 | 732 | 1,407 | 2,958 |
|  | SpC-NB |  | 785 | 1,374 | 2,771 |
|  | Improvement | 5.0x | $\overline{3.3} \overline{\mathbf{x}}$ | $\overline{3} . \overline{7} \mathbf{x}$ | $\overline{2.7 x}{ }^{-}$ |
| $Z=9$ | Dense3D | 981,686 | 1,818 | 3,863 | 6,759 |
|  | $\overline{\mathrm{Sp}} \overline{\mathrm{C}}-\overline{\mathrm{B}}$ - |  | $\overline{9} 0 \overline{2}$ | 1,8 $\overline{2} 0$ | 3,723 ${ }^{-}$ |
|  | SpC-RB | 250,387 | 811 | 1,626 | 2,718 |
|  | SpC-NB |  | 750 | 1,395 | 2,872 |
|  | Improvement | - - $\overline{3.92} \overline{\mathbf{x}}$ | $\overline{2} . \overline{4} 2 \bar{x}$ | $2 . \overline{7} \overline{7} \mathbf{x}$ | $2 . \overline{35 x}{ }^{-}$ |

Table 2 summarizes the runs on 900 processors. A value in the table represents the geometric average of runs along all matrices in the dataset for the respective metric, method, $K$, and $Z$ values. The metrics we consider are maximum receive volume and SDDMM running time. Since the different sparse communication methods within SpComm3D perform the same communication, all of them share the same max receive volume. We report volume values normalized with respect to $K$. We also report the improvement of $\mathrm{SpC}-\mathrm{NB}$ with respect to Dense3D for each different $Z$ value.

As seen in the table, SpComm3D improves the maximum receive volume by 4.0 x to 6.5 x , depending on $Z$. This reduction directly impacts the actual runtime, with improvements from 2.35 x up to 4.7 x , depending on $K$ and $Z$. Compared against each other, $\mathrm{SpC}-\mathrm{BB}$, $\mathrm{SpC}-\mathrm{RB}$, and $\mathrm{SpC-NB}$ show expected behaviour. SpC -NB performs best overall, followed by $\mathrm{SpC}-\mathrm{RB}$. The gap between $\mathrm{SpC}-\mathrm{BB}$ and the other two methods is clear, especially when $K$ gets higher, whereas $\mathrm{SpC-NB}$ and $\mathrm{SpC}-\mathrm{RB}$ are comparable.

Figure 9 shows a breakdown of the SDDMM running times of SpC-NB on 1800 processors with different $K$ and $Z$ values. The figure shows that PreComm phase dominates the running time. The Compute phase's share increase as $K$ increases, whereas the share of PostComm phase increase as $Z$ increases.

## 8 CONCLUSIONS

SpComm3D provides a high-level and flexible environment to run different types of 2D and 3D sparse kernels utilizing sparse communication. With extensive experimental evaluations on up to 1800 processors, we showed that SpComm3D has superior scalability compared to the state-of-the-art sparsity-agnostic SDDMM and SpMM algorithms. The communication and memory required to


Figure 9: Runtime breakdown of SDDMM with SpC-NB of all matrices on 1800 processors (averaged with geometric mean) categorized according to $K$ and $Z$.
run the same algorithm are reduced by an average of 5.0x, leading to saving up to 10 TB of bandwidth and RAM on 1800 processors, by effectively utilizing sparse communication. The overall runtime of SDDMM is also reduced by 5.0 x on average. The low memory and communication overheads of SpComm3D compared to the sparsityagnostic counterparts make it a better candidate for performing large-scale sparse computations especially with very large sparse matrices.

## REFERENCES

[1] Nabil Abubaker, Orhun Caglayan, M Ozan Karsavuran, and Cevdet Aykanat. 2023. Minimizing staleness and communication overhead in distributed SGD for collaborative filtering. IEEE Trans. Comput. (2023).
[2] Seher Acer, Oguz Selvitopi, and Cevdet Aykanat. 2016. Improving performance of sparse matrix dense matrix multiplication on large-scale parallel systems. Parallel Comput. 59 (2016), 71-96.
[3] Ramesh C Agarwal, Susanne M Balle, Fred G Gustavson, Mahesh Joshi, and Prasad Palkar. 1995. A three-dimensional approach to parallel matrix multiplication. IBM Journal of Research and Development 39, 5 (1995), 575-582.
[4] Kadir Akbudak and Cevdet Aykanat. 2014. Parallel Sparse Matrix-Matrix Multiplication Library. Technical Report. Technical report BU-CE-1402, Computer Engineering Department, Bilkent University.
[5] Kadir Akbudak and Cevdet Aykanat. 2014. Simultaneous input and output matrix partitioning for outer-product-parallel sparse matrix-matrix multiplication. SIAM fournal on Scientific Computing 36, 5 (2014), C568-C590.
[6] Kadir Akbudak, Oguz Selvitopi, and Cevdet Aykanat. 2018. Partitioning models for scaling parallel sparse matrix-matrix multiplication. ACM Transactions on Parallel Computing (TOPC) 4, 3 (2018), 1-34.
[7] Ariful Azad, Grey Ballard, Aydin Buluc, James Demmel, Laura Grigori, Oded Schwartz, Sivan Toledo, and Samuel Williams. 2016. Exploiting multiple levels of parallelism in sparse matrix-matrix multiplication. SIAM fournal on Scientific Computing 38, 6 (2016), C624-C651.
[8] Satish Balay, Shrirang Abhyankar, Mark F. Adams, Steven Benson, Jed Brown, Peter Brune, Kris Buschelman, Emil Constantinescu, Lisandro Dalcin, Alp Dener, Victor Eijkhout, Jacob Faibussowitsch, William D. Gropp, Václav Hapla, Tobin Isaac, Pierre Jolivet, Dmitry Karpeev, Dinesh Kaushik, Matthew G. Knepley, Fande Kong, Scott Kruger, Dave A. May, Lois Curfman McInnes, Richard Tran Mills, Lawrence Mitchell, Todd Munson, Jose E. Roman, Karl Rupp, Patrick Sanan, Jason Sarich, Barry F. Smith, Stefano Zampini, Hong Zhang, Hong Zhang, and Junchao Zhang. 2024. PETSc/TAO Users Manual. Technical Report ANL-21/39Revision 3.21. Argonne National Laboratory. https://doi.org/10.2172/2205494
[9] Grey Ballard, Aydin Buluc, James Demmel, Laura Grigori, Benjamin Lipshitz, Oded Schwartz, and Sivan Toledo. 2013. Communication optimal parallel multiplication of sparse random matrices. In Proceedings of the twenty-fifth annual ACM symposium on Parallelism in algorithms and architectures. 222-231.
[10] Maciej Besta and Torsten Hoefler. 2024. Parallel and distributed graph neural networks: An in-depth concurrency analysis. IEEE Transactions on Pattern Analysis and Machine Intelligence (2024).
[11] Vivek Bharadwaj, Aydın Buluç, and James Demmel. 2022. Distributed-memory sparse kernels for machine learning. In 2022 IEEE International Parallel and

Distributed Processing Symposium (IPDPS). IEEE, 47-58.
[12] Aydin Buluc and John R Gilbert. 2008. Challenges and advances in parallel sparse matrix-matrix multiplication. In 2008 37th International Conference on Parallel Processing. IEEE, 503-510.
[13] Aydin Buluç and John R Gilbert. 2012. Parallel sparse matrix-matrix multiplication and indexing: Implementation and experiments. SIAM Journal on Scientific Computing 34, 4 (2012), C170-C191.
[14] Lynn Elliot Cannon. 1969. A cellular computer to implement the Kalman filter algorithm. Montana State University.
[15] Ümit V Çatalyürek, Cevdet Aykanat, and Bora Uçar. 2010. On two-dimensional sparse matrix partitioning: Models, methods, and a recipe. SIAM fournal on Scientific Computing 32, 2 (2010), 656-683.
[16] Timothy A. Davis and Yifan Hu. 2011. The university of Florida sparse matrix collection. ACM Trans. Math. Softw. 38, 1, Article 1 (dec 2011), 25 pages. https: //doi.org/10.1145/2049662.2049663
[17] Gunduz Vehbi Demirci and Hakan Ferhatosmanoglu. 2021. Partitioning sparse deep neural networks for scalable training and inference. In Proceedings of the ACM International Conference on Supercomputing. 254-265
[18] Roger G Grimes, John G Lewis, and Horst D Simon. 1994. A shifted block Lanczos algorithm for solving sparse symmetric generalized eigenproblems. SIAM 7. Matrix Anal. Appl. 15, 1 (1994), 228-272.
[19] Ramakrishnan Kannan, Grey Ballard, and Haesun Park. 2017. MPI-FAUN: An MPI-based framework for alternating-updating nonnegative matrix factorization. IEEE Transactions on Knowledge and Data Engineering 30, 3 (2017), 544-558.
[20] Kamer Kaya, Bora Uçar, and Ümit V. Çatalyürek. 2014. Analysis of Partitioning Models and Metrics in Parallel Sparse Matrix-Vector Multiplication. In Parallel Processing and Applied Mathematics (PPAM, Sep. 2013) (Lecture Notes in Computer Science), Roman Wyrzykowski, Jack Dongarra, Konrad Karczewski, and Jerzy Waśniewski (Eds.). Springer Berlin Heidelberg, Warsaw, Poland, 174-184.
[21] Oguz Kaya, Ramakrishnan Kannan, and Grey Ballard. 2018. Partitioning and communication strategies for sparse non-negative matrix factorization. In Pro ceedings of the 47th International Conference on Parallel Processing. 1-10.
[22] Enver Kayaaslan, Cevdet Aykanat, and Bora Uçar. 2018. 1.5 D parallel sparse matrix-vector multiply. SIAM fournal on Scientific Computing 40, 1 (2018), C25C46

23] Penporn Koanantakool, Ariful Azad, Aydin Buluç, Dmitriy Morozov, Sang-Yun Oh, Leonid Oliker, and Katherine Yelick. 2016. Communication-avoiding parallel sparse-dense matrix-matrix multiplication. In 2016 IEEE International Parallel and Distributed Processing Symposium (IPDPS). IEEE, 842-853.
[24] Yehuda Koren, Robert Bell, and Chris Volinsky. 2009. Matrix factorization techniques for recommender systems. Computer 42, 8 (2009), 30-37.
[25] Grzegorz Kwasniewski, Marko Kabić, Maciej Besta, Joost VandeVondele, Raffaele Solcà, and Torsten Hoefler. 2019. Red-blue pebbling revisited: near optimal parallel matrix-matrix multiplication. In Proceedings of the International Conference for High Performance Computing, Networking, Storage and Analysis. 1-22.
[26] Ujjaini Mukhopadhyay. 2023. Sparsity-aware communication for distributed graph neural network training. Master's thesis. EECS Department, University of California, Berkeley. http://www2.eecs.berkeley.edu/Pubs/TechRpts/2023/EECS-2023-253.html
[27] Dianne P O'Leary. 1980. The block conjugate gradient algorithm and related methods. Linear algebra and its applications 29 (1980), 293-322.
[28] Miloud Sadkane. 1993. A block Arnoldi-Chebyshev method for computing the leading eigenpairs of large sparse unsymmetric matrices. Numerische mathematik 64 (1993), 181-193.
[29] Oguz Selvitopi, Benjamin Brock, Israt Nisa, Alok Tripathy, Katherine Yelick, and Aydın Buluç. 2021. Distributed-memory parallel algorithms for sparse times tall-skinny-dense matrix multiplication. In Proceedings of the ACM International Conference on Supercomputing (Virtual Event, USA) (ICS '21). Association for Computing Machinery, New York, NY, USA, 431-442. https://doi.org/10.1145/ 3447818.3461472
[30] Edgar Solomonik and James Demmel. 2011. Communication-optimal parallel 2.5 D matrix multiplication and LU factorization algorithms. In European Conference on Parallel Processing. Springer, 90-109.
[31] Alok Tripathy, Katherine Yelick, and Aydın Buluç. 2020. Reducing Communication in Graph Neural Network Training. In SC20: International Conference for High Performance Computing, Networking, Storage and Analysis. 1-14. https://doi.org/10.1109/SC41405.2020.00074
[32] Robert A Van De Geijn and Jerrell Watts. 1997. SUMMA: Scalable universal matrix multiplication algorithm. Concurrency: Practice and Experience 9, 4 (1997), 255-274.


[^0]:    ${ }^{1}$ https://sparse.tamu.edu/

